Biological Regulatory Networks:

Logical Description

and Model Checking

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Programme ÉPIGÉNOMIQUE, Genopole®

Special thanks to J. Guespin, J-P. Comet & the Observability group





Menu

- 1. Modelling biological regulatory networks
- 2. Formal framework for biological regulatory networks
- 3. Temporal logic and Model Checking
- 4. Computer aided elaboration of formal models
- 5. Example: mucus production in *Pseudomonas aeruginosa*

Molecular Biology & Causality

Heaviness of "causality networks"

Causality loops

Counter-intuitive resulting behaviours

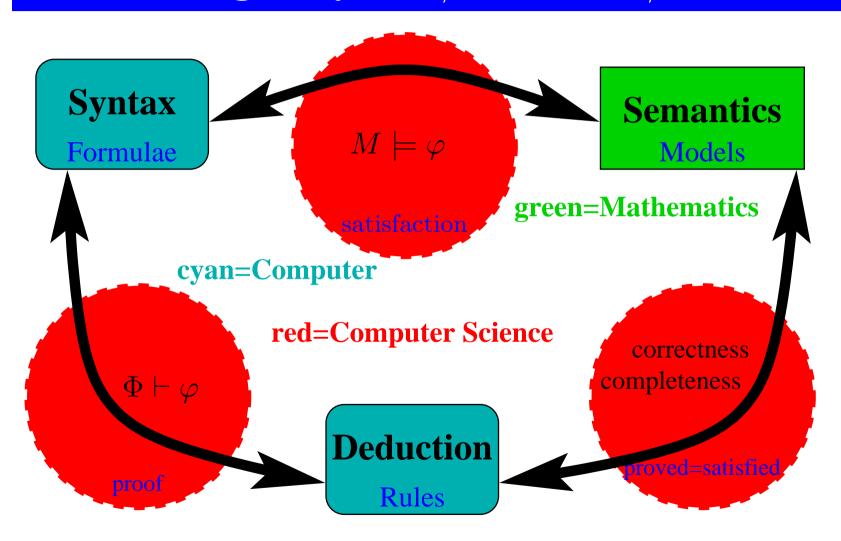
Predicting dynamics from models

Computer aided modelling methodologies

- \bullet quantitative approaches \rightarrow e.g. differential equations
- ullet qualitative approaches \to logic & computer science can help
- mixed approaches \rightarrow cf. Marcelline Kaufman

Biological questions are often of qualitative nature

Formal Logic: syntax/semantics/deduction



Regulatory Networks

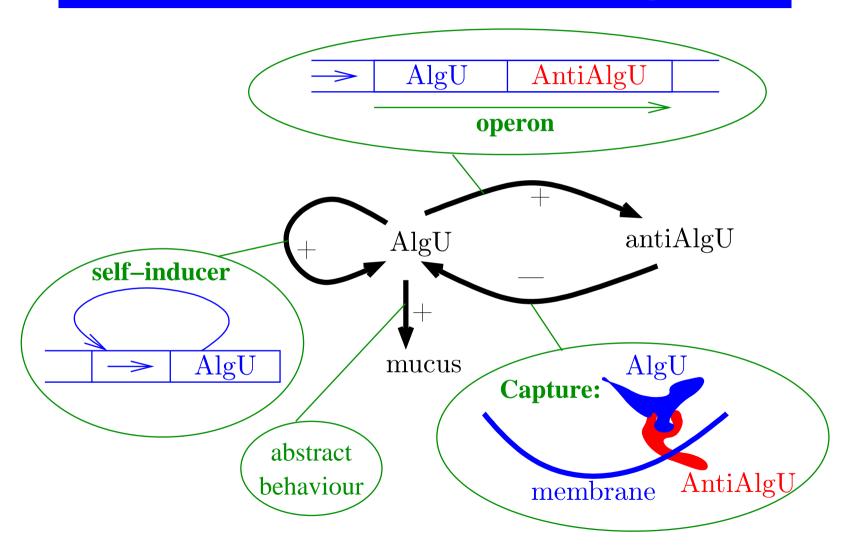
To model direct or indirect regulations between biological objects (e.g. gene, macromolecule, signal, ...)

Direct: transcription factor, operon, repressor, ...

Indirect: cascade of events, capture of macromolecules, . . .

- x induces $y: x \xrightarrow{+} y$ x inhibits $y: x \xrightarrow{-} y$

Mucus Production in P. aeruginosa

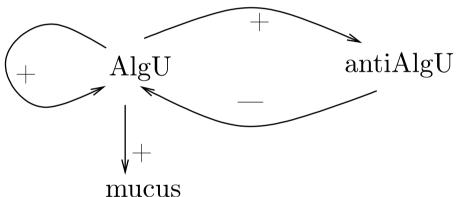


Static Graph & Dynamic Behaviour

Difficulty to predict the result of combined regulations

Difficulty to measure the strength of a given regulation

Example of "competitor" circuits



Positive v.s. Negative circuits

Even v.s. Odd number of "—" signs

Multistationarity v.s. Homeostasy

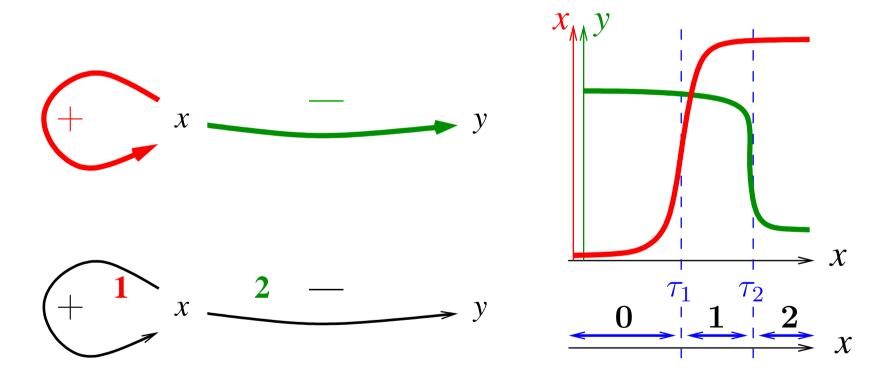
René Thomas, Snoussi, ..., Soulé

Functional circuits "pilot" the behaviour

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Multivalued Regulatory Graphs



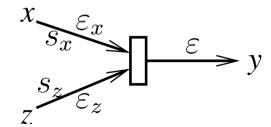
Definition of Regulatory Graphs

A labelled directed graph $(\mathcal{V}, \mathcal{E})$

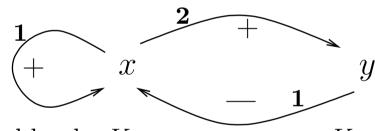
- each node of \mathcal{V} is a variable x with a boundary $b_x \in I\!\!N$, less or equal to the out-degree of x.
- each edge $x \to y$ of \mathcal{E} is labelled by $\varepsilon \in \{+, -\}$ and by $s \in [0 \cdots b_x]$.

Variant: bipartite graph

- complexation of two proteins
- inhibition of a regulation
- external conditions...



Regulatory Networks (R. Thomas)



 K_y Basal level : K_x $K_{y,x}$

 $x \text{ helps} : K_{x,x}$

Absent y helps : $K_{x,\overline{y}}$

Both : $K_{x,x\overline{y}}$

(x,y)	\underline{Image}
(0,0)	$(K_{x,\overline{y}},K_y)$
(0,1)	(K_x, K_y)
(1,0)	$(K_{x,x\overline{y}},K_y)$
(1,1)	$(K_{x,x},K_y)$
(2,0)	$(K_{x,x\overline{y}},K_{y,x})$
(2,1)	$(K_{x,x},K_{y,x})$

Resources in a Regulatory Network

States:

 $\eta: \mathcal{V} \to I\!\!N \ (\approx \text{vector of integers})$ $\eta(x) = abstract \ concentration \ level \ of \ x$

Variant: singular states (values can be the thresholds τ_1, τ_2, \ldots)

Resources:

For each $x \xrightarrow{+,s} y$, x is a resource of y iff $\eta(x) > s$

For each $x \xrightarrow{-,s} y$, x is a resource of y iff $\eta(x) \leq s$

Parameters:

Partial function $K: \mathcal{V} \times \mathcal{P}(\mathcal{V}) \to I\!\!N$

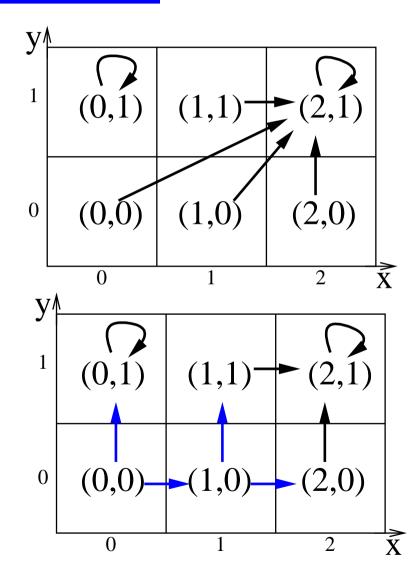
Image:

Vector of the $K(y,\omega)$ where ω is the set of resources of y

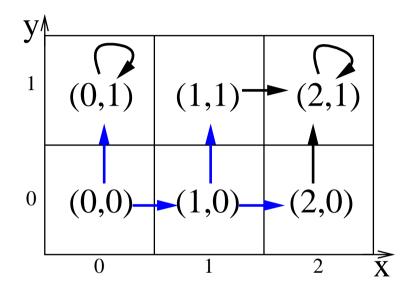
State Graphs

(x,y)	\underline{Image}
(0,0)	$(K_{x,\overline{y}},K_{y})=(2,1)$
(0,1)	$(K_x, K_y) = (0,1)$
(1,0)	$(K_{x,x\overline{y}},K_y)=(2,1)$
(1,1)	$(K_{x,x}, K_y) = (2,1)$
(2,0)	$(K_{x,x\overline{y}},K_{y,x})=(2,1)$
(2,1)	$(K_{x,x}, K_{y,x}) = (2,1)$

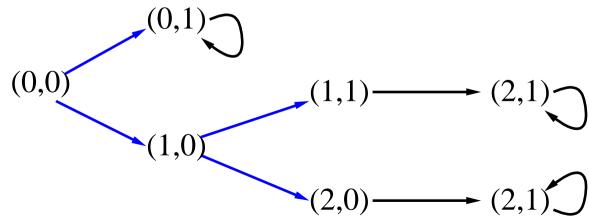
"desynchronization" \longrightarrow by units of Manhattan distance



Time has a tree structure

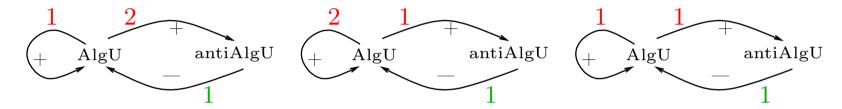


From an initial state:



Parameters & thresholds: often unknown

Thresholds for AlgU in *P.aeruginosa* are unknown:



and parameters are unknown:

$$3^4 \times 2^2$$

$$2^4 \times 2^2$$

712 possible models

 $3^4 \times 2^2$

Some criteria exist to reduce the number of models,

but formal logic is needed to go further automatically

Note: some models are observably equivalent

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CTL = Computation Tree Logic

Atoms = comparaisons : (x=2) (y>0) ...

Logical connectives: $(\varphi_1 \land \varphi_2) \quad (\varphi_1 \implies \varphi_2)$

Temporal connectives: made of 2 characters

first character

 $A = \text{for All path choices} \mid X = \text{neXt state}$

E =there **E**xist a choice

second character

F =for some Future state

G =for all future states (Globally)

 $U = \mathbf{U}$ ntil

AX(y=1): the concentration level of y belongs to the interval 1 in all states directly following the considered initial state.

EG(x=0): there exists at least one path from the considered initial state where x always belongs to its lower interval.

Temporal Connectives of CTL

neXt state:

 $EX\varphi: \varphi$ can be satisfied in a next state

 $AX\varphi: \varphi$ is always satisfied in the next states

eventually in the Future:

 $EF\varphi: \varphi$ can be satisfied in the future

 $AF\varphi: \varphi$ will be satisfied at some state in the future

Globally:

 $EG\varphi: \varphi$ can be an invariant in the future

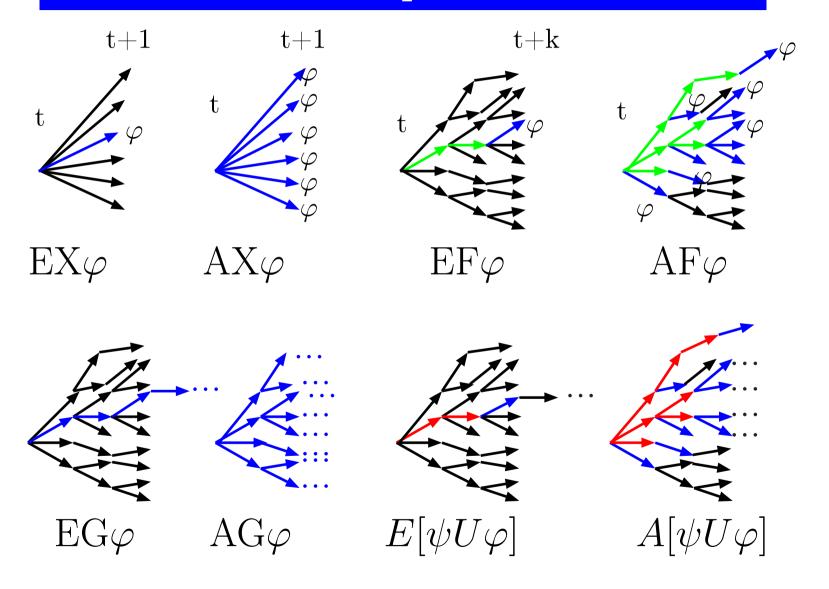
 $AG\varphi: \varphi$ is necessarily an invariant in the future

Until:

 $E[\psi U\varphi]$: there exist a path where ψ is satisfied until a state where φ is satisfied

 $A[\psi U\varphi]$: ψ is always satisfied until some state where φ is satisfied

Semantics of Temporal Connectives

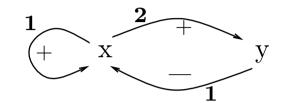


CTL to encode Biological Properties

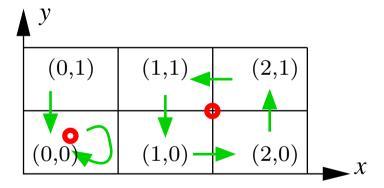
Common properties:

"functionality" of a sub-graph

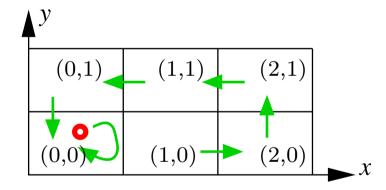
Special role of "feedback loops"



- positive: multistationnarity (even number of)
- negative: homeostasy (odd number of)



Characteristic properties:



$$(x=2) \Longrightarrow AG(\neg(x=0))$$

and
$$(x=0) \Longrightarrow AG(\neg(x=2))$$

They express "the positive feedback loop is functional" (satisfaction of these formulae relies on the parameters $K_{...}$)

$\overline{\textbf{Theoretical Models}} \leftrightarrow \overline{\textbf{Experiments}}$

CTL formulae are satisfied (or refuted) w.r.t. a set of paths from a given initial state

- They can be tested against the possible paths of the theoretical models $(M \models_{\eta} \varphi)$
- They can be tested against the biological experiments $(Biological_Object \models_{experiment} \varphi)$

CTL formulae link theoretical models and biological objects together

Model Checking

Computes all the states of a theoretical model which satisfy a given formula: $\{ \eta \mid M \models_{\eta} \varphi \}$.

Idea 1: work on the state graph instead of the path trees.

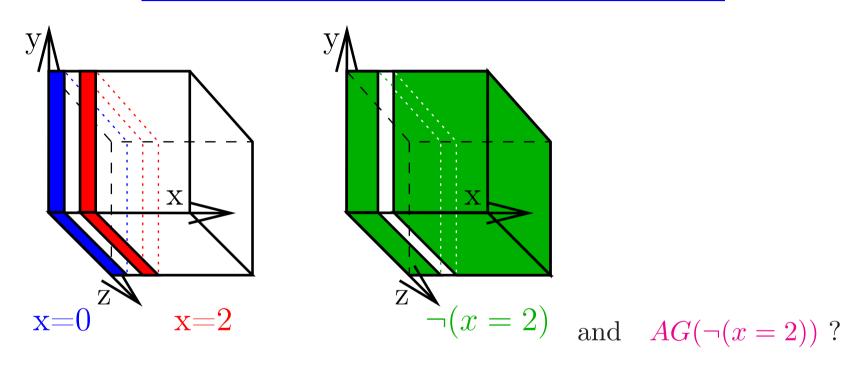
Idea 2: check first the atoms of φ and then check the connectives of φ with a bottom-up computation strategy.

Idea 3: (computational optimization) group some cases together using BDDs (Binary Decision Diagrams).

Example:
$$(x=0) \implies AG(\neg(x=2))$$

Obsession: travel the state graph as less as possible





... one should **travel** <u>all</u> the paths from any green box and check if successive boxes are green: *too many boxes to visit*.

Trick: $AG(\neg(x=2))$ is equivalent to $\neg EF(x=2)$ start from the red boxes and follow the transitions backward.

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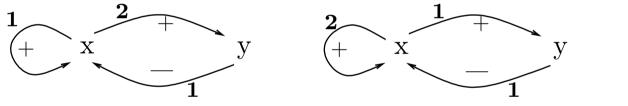
Computer Aided Elaboration of Models

From biological knowledge and/or biological hypotheses, it comes:

• properties:

"Without stimulus, if gene x has its basal expression level, then it remains at this level."

• model schemas:



Formal logic and formal models allow us to:

- verify hypotheses and check consistency
- elaborate more precise models incrementally
- suggest new biological experiments to efficiently reduce the number of potential models

The Two Questions



1. Is it possible that Φ and \mathcal{M} ?

Consistency of knowledge and hypotheses. Means to select models belonging to the schemas that satisfy Φ .

$$(\exists ? M \in \mathcal{M} \mid M \models \varphi)$$

2. If so, is it true in vivo that Φ and \mathcal{M} ?

Compatibility of one of the selected models with the biological object. Require to propose experiments to **validate** (or refute) the selected model(s).

 \rightarrow Computer aided *proofs* and *validations*

Question 1 = Consistency

- 1. Draw all the sensible regulatory graphs with all the sensible threshold allocations. It defines \mathcal{M} .
- 2. Express in CTL the known behavioural properties as well as the considered biological hypotheses. It defines Φ .
- 3. Automatically generate all the possible regulatory networks derived from \mathcal{M} according to all possible parameters $K_{...}$. Our software plateform SMBioNet handles this automatically.
- 4. Check each of these models against Φ . SMBioNet uses model checking to perform this step.
- 5. If no model survive to the previous step, then reconsider the hypotheses and perhaps extend model schemas...
- 6. If at least one model survives, then the biological hypotheses are consistent. Possible parameters $K_{...}$ have been indirectly established. Now Question 2 has to be addressed.

Question 2 = Validation

- 1. Among all possible formulae, some are "observable" i.e., they express a possible result of a possible biological experiment. Let *Obs* be the set of all observable formulae.
- 2. Let Λ be the set of theorems of Φ and \mathcal{M} . $\Lambda \cap Obs$ is the set of experiments able to validate the survivors of Question 1. Unfortunately it is infinite in general.
- 3. Testing frameworks from computer science aim at selecting a finite subsets of these observable formulae, which maximize the chance to refute the survivors.
- 4. These subsets are often too big but in some cases, these testing frameworks can be applied to regulatory networks.

 It has been the case of the mucus production of *P.aeruginosa*.

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Mutation, Epigenesis, Adaptation

Terminology about phenotype modification:

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genetic modification: inheritable and not reversible (mutation)epigenetic modification: inheritable and reversibleadaptation: not inheritable and reversible
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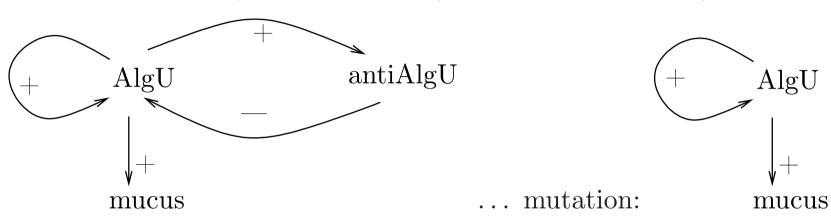
The biological questions (Janine Guespin): are mucus production and/or cytotoxicity in *Pseudomonas aeruginosa* due to an epigenetic switch ?

[—— cystic fibrosis]

Mucus production in P. aeruginosa

Pseudomonas aeruginosa:

(J.Guespin, M.Kaufman)

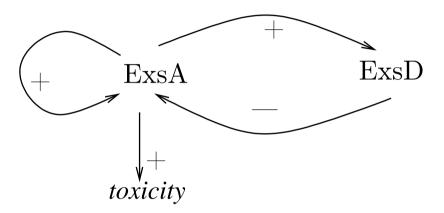


Epigenetic Hypothesis (without mutation) =

- \rightarrow The positive feedback circuit is functional, with a mucoid stable state and another non mucoid stable state.
- → An external signal (in the cystic fibrosis' lungs) could switch AlgU from its lower stable state to the higher one.
- \rightarrow The mutation could be favored later because the inhibitor complex is toxic for the bacteria. \Longrightarrow New possible therapy.

Cytotoxicity in *P. aeruginosa*

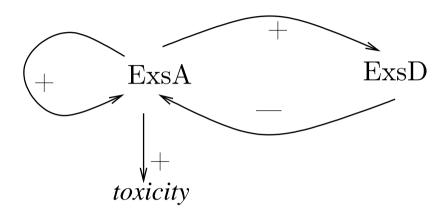
(Janine Guespin)



Epigenetic hypothesis =

- \rightarrow The positive feedback circuit is functional, with a cytotoxic stable state and the other one is not cytotoxic.
- → An external signal (in the cystic fibrosis' lungs) could switch ExsA from its lower stable state to the higher one.

Consistency of the Hypothesis



One CTL formula for each stable state:

$$(ExsA = 2) \Longrightarrow AXAF(ExsA = 2)$$

$$(ExsA = 0) \Longrightarrow AG(\neg(ExsA = 2))$$

Question 1, consistency: proved by Model Checking

 \rightarrow 10 models among the 712 models are extracted by SMBioNet

Question 2: and in vivo? ...

Validation of the epigenetic hypothesis

Question $2 = \text{to validate bistationnarity } in \ vivo$

Non cytotoxic state: $(ExsA = 0) \Longrightarrow AG(\neg(ExsA = 2))$

P. aeruginosa, with a basal level for ExsA does not become spontaneously cytotoxic: actually validated

Cytotoxic state:
$$(ExsA = 2) \Longrightarrow AXAF(ExsA = 2)$$

Experimental limitation:

ExsA can be saturated but it cannot be measured.

Experiment:

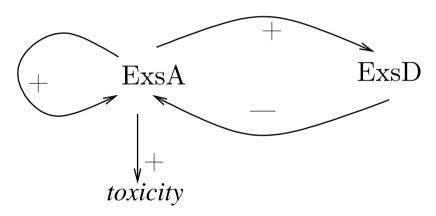
to pulse ExsA and then to test if toxin production remain.

 \iff to verify a hysteresis)

This experiment can be generated automatically

To test $(ExsA=2) \Longrightarrow AXAF(ExsA=2)$

ExsA = 2 cannot be directly verified but toxicity = 1 can be verified.



Lemma: $AXAF(ExsA = 2) \iff AXAF(toxicity = 1)$ (... formal proof by computer ...)

$$\rightarrow$$
 To test: (ExsA = 2) $\Longrightarrow AXAF(toxicity = 1)$

$(ExsA = 2) \Longrightarrow AXAF(toxicity = 1)$

Karl Popper:

$A \Longrightarrow B$	true	false
true	true	false
false	true	true

to validate = to try to refute $thus \ A = false \ is \ useless$ experiments must begin with a pulse

The pulse forces the bacteria to reach the initial state ExsA = 2. If the state were not directly controlable we had to prove lemmas:

$$(ExsA = 2) \iff (something\ reachable)$$

General form of a test:

 $(something \ \underline{reachable}) \Longrightarrow (something \ \underline{observable})$

Concluding Slogans

- Behavioural properties (Φ) are as much important as models (\mathcal{M}) for the modelling activity
- Modelling is significant only with respect to the considered experimental reachability and observability (Obs)
- The bigger is the risk of *refutation*, the better are the "surviving" models (Popper), thus models should be "simple" with few non observable parameters (Occam)

Formal methods (syntax/semantics/proofs) facilitate abstraction and consequently they simplify models

- They ensure *consistency* of the modelling activity
- They allow us to perform computer aided *validations* of models
- They take benefit of 30 years of researches in computer sciences